

# Problem Formulation

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$$\begin{aligned} f(T_{min_i}, r_i) &= \sum q(v_{min})|p(r_i, v_{min})| \\ &= q(v_1)|p(r_i, v_1)| + q(v_2)|p(r_i, v_2)| + \dots + q(v_{V-1})|p(r_i, v_{V-1})| \end{aligned} \tag{1}$$

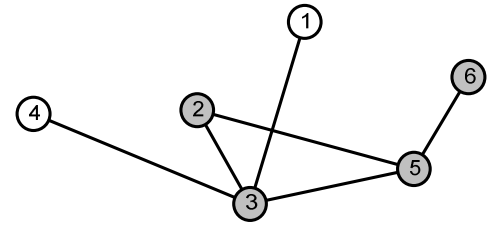
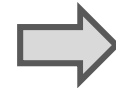
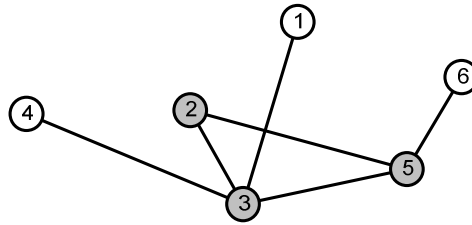
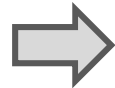
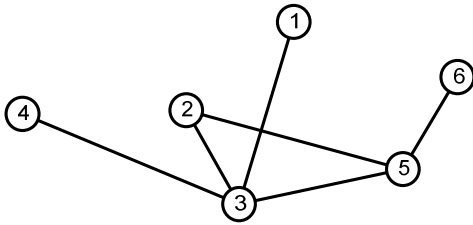


$$\begin{aligned} f'(T_{min_i}, r_i) &= \sum |p(r_i, v_{min})| \\ &= |p(r_i, v_1)| + |p(r_i, v_2)| + \dots + |p(r_i, v_{V-1})| \end{aligned} \tag{2}$$

**Assumption 1** — The local best solution for equation (1) provides the best overall solution.

**Assumption 2** — The derived sub-graph forms a spanning tree structure.

# Algorithm Optimization



Vertex #	Degree	Weight
4	1	4
2	2	2
3	4	3
1	1	1
5	3	5
6	1	6

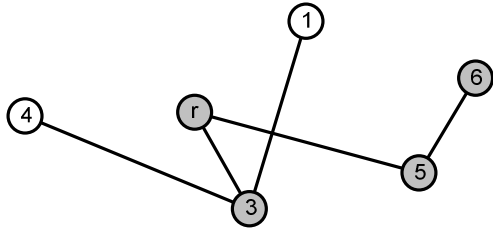
Vertex #	Degree	Weight
4	1	4
1	1	1
6	1	6

Vertex #	Degree	Weight
2	2	2
3	4	3
5	3	5

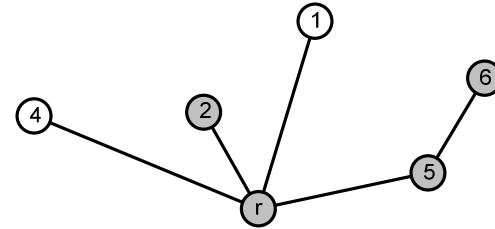
Vertex #	Degree	Weight
4	1	4
1	1	1

Vertex #	Degree	Weight
2	2	2
3	4	3
5	3	5
6	1	6

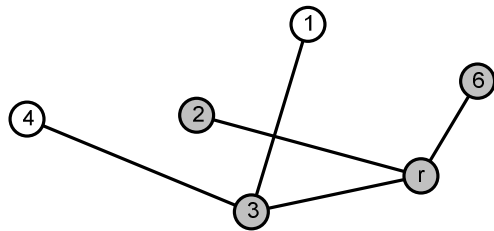
# Algorithm Design



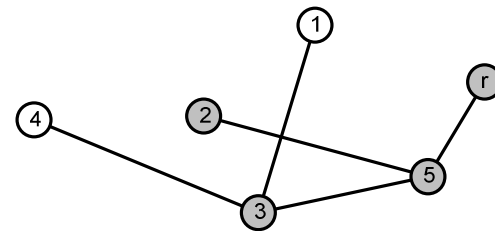
$$\begin{aligned}
 f(T, r) &= \sum q(v)|p(r, v)| \\
 &= 4 \times 2 + 3 \times 1 + 1 \times 2 + 5 \times 1 + 6 \times 2 \\
 &= 30
 \end{aligned}$$



$$\begin{aligned}
 f(T, r) &= \sum q(v)|p(r, v)| \\
 &= 4 \times 1 + 2 \times 1 + 1 \times 1 + 5 \times 1 + 6 \times 2 \\
 &= 24
 \end{aligned}$$



$$\begin{aligned}
 f(T, r) &= \sum q(v)|p(r, v)| \\
 &= 4 \times 2 + 2 \times 1 + 1 \times 2 + 3 \times 1 + 6 \times 1 \\
 &= 21 \quad \leftarrow
 \end{aligned}$$



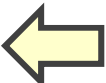
$$\begin{aligned}
 f(T, r) &= \sum q(v)|p(r, v)| \\
 &= 4 \times 3 + 2 \times 2 + 1 \times 3 + 3 \times 2 + 5 \times 1 \\
 &= 32
 \end{aligned}$$

# Pseudo-code

## JADE-BFS-ALGORITHM(Matrix $M$ )

```
1  int  $maxWeight \leftarrow 0$ ;  
2  int  $solutionSum \leftarrow 0$ ;  
3  Matrix  $input \leftarrow$  the input adjacency  $M$ ;  
4  Matrix  $solution \leftarrow$  populate with 0;  
5  VertexNode  $root \leftarrow$  null;  
6  Queue  $candidateSet \leftarrow \emptyset$   
7  for every  $v \in V$  {  
8      if  $v.getWeight > maxWeight$  {  
9           $maxWeight \leftarrow v.getWeight$   
10     }  
11 }  
12 for every  $v \in V$  {  
13     if  $v.getDegree \neq 1$  or  $v.getWeight = mWeight$  {  
14          $candidateSet.push(v)$ ;  
15     }  
16 }
```

```
17 while  $candidateSet \neq \emptyset$  {  
18     Matrix  $localSolution \leftarrow$  populate with 0;  
19     int  $localSum \leftarrow 0$ ;  
20     Queue  $temp \leftarrow \emptyset$ ;  
21     VertexNode  $vroot \leftarrow candidateSet.pop$ ;  
22      $temp.push(vroot)$ ;  
23      $vroot.setHops(0)$ ;  
24     for every  $v \in V$  and  $v \neq vroot$  {  
25          $v.setHops(\infty)$   
26     }  
27     while  $temp \neq \emptyset$  {  
28         VertexNode  $u \leftarrow temp.pop$ ;  
29         for every  $v \in V$  and  $v \neq vroot$  {  
30             if  $input.adjacent(v, u)$  {  
31                 if  $v.getHops = \infty$  {  
32                      $v.setHops(1 + u.getHops)$ ;  
33                      $localSolution.mark(v, u)$ ;  
34                      $temp.push(v)$ ;  
35                 }  
36             }  
37         }  
38     }  
39     for every  $v \in V$  {  
40          $localSum \leftarrow localSum + v.getWeight \cdot v.getHops$ ;  
41     }  
42     if  $localSum < solutionSum$  {  
43          $solutionSum \leftarrow localSum$   
44          $solution \leftarrow localSolution$   
45          $root \leftarrow vroot$   
46     }  
47 }
```



# Complexity Analysis

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## *Space Complexity*

***InputMatrix + VertexNode[V] + OutputMatrix + OutputVertexNode + Queue***

$$\begin{aligned}\text{Space Complexity} &= V \times \frac{V}{2} + V \times 3 + V \times \frac{V}{2} + E \\ &= V \times V + V \times 3 + E \\ &= V^2 + V \times 3 + E\end{aligned}$$

## *Time Complexity*

**Loop 7—11 + Loop 12—15 + Loop 16—47 × Loop 27—38 × Loop 30—37**

$$\begin{aligned}\text{Time Complexity} &= O(V) + O(V) + O(V) \times O(V) \times O(D) \\ &= 2 \times O(V) + O(V) \times O(E) \\ &= O(V \times E)\end{aligned}$$